## Continuous max-flow and global minimization for high dimensional data

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## Interface (classification) problems

Interface problems exists everywhere in science and technology. For imaging and vision, it is somehow classical:

- Mumford-Shal model
- GAC model
- Chan-Vese model

How to solve these interface problems?

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- active count our
- level set
- phase-field
- ...


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- active count our
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Traditional methods are:

- Nonlinear
- Non-convex
- ...


## Max-Flow / Min-Cut



## Max-Flow / Min-Cut


$\left(V_{s}, V_{t}\right)$ is a cut, $w_{i j}=$ cost of cutting edge $(i, j)$ $\operatorname{cost}$ of cut $c\left(V_{s}, V_{t}\right)=\sum_{i \in V_{s}, j \in V_{t}} w_{i j}$

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## Higher dimensional problems

A graph for 2D images:


Figure: Graph used for discrete 2D binary labeling

## Two-phase Min-cut - Discretized setting



Figure: Graph and cut for discrete binary labeling
It is easy to see the cost of a cut $(u(p)=0$ or 1$)$. A minimum cut is to find $u$ for:

$$
\min _{u \in\{0,1\}} \sum_{p \in \mathcal{P}} f_{1}(p)(1-u(p))+f_{2}(p) u(p)+\sum_{p \in \mathcal{P}} \sum_{q \in \mathcal{N}_{p}^{k}} g(p, q)|u(p)-u(q)| .
$$

Capacity:

$$
w_{s, p}=f_{1}(p), w_{t, p}=f_{2}(p), w_{p, q}=g(p, q)
$$

## Two-phase Min-cut - corresponding continuous setting



Figure: Graph used for discrete and continuous binary labeling A "continuous" minimum cut is to solve:
$\min _{u \in\{0,1\}} \int_{\Omega} f_{1}(x)(1-u(x))+f_{2}(x) u(x)+g_{1}(x)\left|D_{1} u(x)\right|+g_{2}(x)\left|D_{2} u(x)\right|$.
Capacity:

$$
w_{s}(x)=f_{1}(x), w_{t}(x)=f_{2}(x), w_{1}(x)=g_{1}(x), \quad w_{2}(x)=g_{2}(x)
$$

## Max-Flow over a graph



Figure: Graph used for discrete binary labeling

## Max-flow formulation

$$
\max _{p_{s}, p_{t}, q} \sum_{v \in \mathcal{V} \backslash\{s, t\}} p_{s}(v)
$$

subject to

$$
\begin{array}{ll}
|q(v, u)| \leq g(v, u), & \forall(v, u) \in \mathcal{V} \times \mathcal{V} \\
0 \leq p_{s}(v) \leq f_{1}(v), & \forall v \in \mathcal{V} \backslash\{s, t\} \\
0 \leq p_{t}(v) \leq f_{2}(v), & \forall v \in \mathcal{V} \backslash\{s, t\}
\end{array}
$$

$$
\left(\sum_{u \in N(v)} q(v, u)\right)-p_{s}(v)+p_{t}(v)=0, \quad \forall v \in \mathcal{V} \backslash\{s, t\}
$$

## Continuous Max-Flow



Figure: Discrete (left) vs. Continuous (right)
Continuous max-flow formulation

$$
\sup _{p_{s}, p_{t}, q} \int_{\Omega} p_{s}(x) d x
$$

subject to ${ }_{\left|q_{1}(x)\right| \leq g_{1}(x) ; ~\left|q_{2}(x)\right| \leq g_{2}(x), \quad \forall x \in \Omega ; ~}^{\text {; }}$

$$
\begin{aligned}
0 \leq p_{s}(x) \leq f_{1}(x), & \forall x \in \Omega ; \\
0 \leq p_{t}(x) \leq f_{2}(x), & \forall x \in \Omega ; \\
\operatorname{div} q(x)-p_{s}(x)+p_{t}(x)=0, & \text { a.e. } x \in \Omega
\end{aligned}
$$

Related: (G. Strang (1983)).

## Continuous Max-Flow: different internal flow capacity



Figure: Discrete (left) vs. Continuous (right)

Continuous max-flow formulation

$$
\sup _{p_{s}, p_{t}, q} \int_{\Omega} p_{s}(x) d x
$$

subject to

$$
\begin{aligned}
|q(x)|=\sqrt{q_{1}^{2}(x)+q_{2}^{2}(x)} \leq g(x), & \forall x \in \Omega \\
0 \leq p_{s}(x) \leq f_{1}(x), & \forall x \in \Omega \\
0 \leq p_{t}(x) \leq f_{2}(x), & \forall x \in \Omega \\
\operatorname{div} q(x)-p_{s}(x)+p_{t}(x)=0, & \text { a.e. } x \in \Omega
\end{aligned}
$$

## Connection: Continuous Max-Flow and Min-Cut

Lagrange multiplier $u$ for flow conservation condition

$$
\operatorname{div} q(x)-p_{s}(x)+p_{t}(x)=0, \quad \text { a.e. } x \in \Omega
$$

yields primal-dual formulation

$$
\sup _{p_{s}, p_{t}, q} \inf _{u} \int_{\Omega} p_{s}+u\left(\operatorname{div} q-p_{s}+p_{t}\right) d x
$$

s.t. $\quad p_{s}(x) \leq f_{1}(x), \quad p_{t}(x) \leq f_{2}(x), \quad|q(x)| \leq g(x)$.

Optimizing for flows $p_{s}, p_{t}, q$ results in:

$$
\min _{u \in[0,1]} \int_{\Omega} f_{1}(x)(1-u(x))+f_{2}(x) u(x) d x+g(x)|\nabla u(x)| d x
$$

This is exactly the same model as the model in CEN (2006). ${ }^{1}$

[^0]
## Three problems

PCLMS or Binary LM (Lie-Lysaker-T.,2005):

$$
\min _{u(x) \in\{0,1\}} \int_{\Omega} f_{1}(1-u)+f_{2} u+g(x)|\nabla u| d x
$$

Convex problem (CEN, (Chan-Esdoglu-Nikolova,2006))

$$
\min _{u(x) \in[0,1]} \int_{\Omega} f_{1}(1-u)+f_{2} u+g(x)|\nabla u| d x .
$$

Graph-cut (Boykov-Kolmogorov, 2001)

$$
\begin{aligned}
& \max _{p_{s}, p_{t}, q} \int_{\Omega} p_{s} d x \text { subject to: } \\
& p_{s}(x) \leq f_{1}(x), p_{t}(x) \leq f_{2}(x),|p(x)| \leq g(x) \\
& \operatorname{div} p(x)-p_{s}(x)+p_{t}(x)=0
\end{aligned}
$$

## Continuous Max-Flow and Min-Cut

## Multiplier-Based Maximal-Flow Algorithm

Augmented lagrangian functional (Glowinski \& Le Tallec, 1989)
$L_{c}\left(p_{s}, p_{t}, q, \lambda\right):=\int_{\Omega} p_{s} d x+\lambda\left(\operatorname{div} q-p_{s}+p_{t}\right)-\frac{c}{2}\left|\operatorname{div} q-p_{s}+p_{t}\right|^{2} d x$.
minmax subject to:
$p_{s}(x) \leq f_{1}(x), \quad p_{t}(x) \leq f_{2}(x), \quad|q(x)| \leq g(x)$
ADMM algorithm: For $\mathrm{k}=1, \ldots$ until convergence, solve

$$
\begin{aligned}
q^{k+1} & :=\arg \max _{\|q\|_{\infty} \leq \alpha} L_{c}\left(p_{s}^{k}, p_{t}^{k}, q, \lambda^{k}\right) \\
p_{s}^{k+1} & :=\arg \max _{p_{s}(x) \leq f_{1}(x)} L_{c}\left(p_{s}, p_{t}^{k}, q^{k+1}, \lambda^{k}\right) \\
p_{t}^{k+1} & :=\arg \max _{p_{t}(x) \leq f_{2}(x)} L_{c}\left(p_{s}^{k+1}, p_{t}, q^{k+1}, \lambda^{k}\right) \\
\lambda^{k+1} & =\lambda^{k}-c\left(\operatorname{div} q^{k+1}-p_{s}^{k+1}+p_{t}^{k+1}\right)
\end{aligned}
$$

## Metrication error, Parallel, GPU, ...





Experiment of mean-curvature driven 3D surface evolution (volume size: 150X150X150 voxels). (a) The radius plot of the 3D ball evolution driven by its mean-curvature flow, which is computed by the proposed continuous max-flow algorithm; its function is theoretically $r(t)=\sqrt{C-2 t}$. (b) The computed 3D ball at one discrete time frame, which fits a perfect 3D ball shape. This is in contrast to (c), the computation result by graph cut [15] with a 3D 26 -connected graph. The computation time of the continuous max-flow algorithm for each discrete time evolution is around 1 sec ., which is faster than the graph cut method ( 120 sec .)

Ref: Y. Yuan, E. Ukwatta, X. Tai, A. Fenster, and C. Schnorr. A fast global optimization-based approach to evolving contours with generic shape prior. Technical report, also UCLA Tech. Report CAM 12-38, 2012.

## Metrication error, Parallel, GPU, ...

- Fully parallel, easy GPU implementation.
- linear grow of computational cost (per iteration): 2D, 3D, ...


## Multiphase Approaches

Multiphase Approaches

## Multiphase problems - Approach I

We need to identify $n$ characteristic functions

$$
\begin{aligned}
& \psi_{i}(x), i=1,2 \cdots n: \\
& \psi_{i}(x) \in\{0,1\}, \quad \sum_{i=1}^{n} \psi_{i}(x)=1
\end{aligned}
$$


$u_{i}(x) \in\left\{a_{0} 1\right\}, \sum u_{1}(x)=1$


Ref: Zach-et-al (2008), Lelmman-et-al(2009,2010, 2013), Bae-Yuan-T. (IJCV 2009), Bae-et-al (SSVM2009, 2012), Yuan-et-al (ECCV2010, (VPR2010)

## Multiphase problems - Approach II

Each point $x \in \Omega$ is labelled by a vector function:

$$
u(x)=\left(u_{1}(x), u_{2}(x), \cdots u_{d}(x)\right)
$$

Ref: Vese-Chan (2002), Bae-T. (JMIV2013), Lie-et-al (IEEE TIP 2006), Bae-et-al(2012,2013), Liu-T.-Leung(EMMCVPR2013), Nieuwenhuis-Toppe-Cremers (IJCV2013).

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- Multiphase: Total number of phases $n=2^{d}$. (Chan-Vese)


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- Multiphase: Total number of phases $n=2^{d}$. (Chan-Vese)

- More than binary labels: Total number of phases $n=B^{d}$.

$$
u_{i}(x) \in\{0,1,2, \cdots B\} .
$$

Ref: Vese-Chan (2002), Bae-T. (JMIV2013), Lie-et-al (IEEE TIP 2006), Bae-et-al(2012,2013), Liu-T.-Leung(EMMCVPR2013), Nieuwenhuis-Toppe-Cremers (IJCV2013).

## Multiphase problems - Approach III

Each point $x \in \Omega$ is labelled by

$$
u(x)=i, \quad i=1,2, \cdots n
$$

- One label function is enough for any $n$ phases.
- More generall


$$
u(x)=\ell_{i}, \quad i=1,2, \cdots n
$$

Ref: Lie-Lysaker-T. (2006), Ishikawa(2003), Darbon-Sigelle (2006), Pock-et-al(2008), Bae-T. (SSVM2009), Brown-Chan-Bresson(2011), Bae-Yuan-T.-Boykov $(2013,2014)$.

## Application to machine learning

We would like to classify a high dimensional data into multi-classes.


## The model

- Graph: $G=(V, E)$, each data is one vertices in $V, E$ are the connections.
- the weight on $E$ :

$$
\begin{equation*}
w(x, y)=e^{-\frac{d(x, y)^{2}}{\sigma^{2}}} \tag{1}
\end{equation*}
$$

$d(x, y)$ : distance measure.
Another choice:

$$
\begin{equation*}
w(x, y)=e^{-\frac{d(x, y)^{2}}{\sqrt{\tau(x) \tau(y)}}} \tag{2}
\end{equation*}
$$

$\tau(x)=d(x, z)^{2}$, and $z$ is the $M^{t h}$ closest vertex to vertex $x$.

## Nonlocal TV on a graph

Gradient operator $\nabla: \mathcal{V} \rightarrow \mathcal{E}$ is:

$$
\begin{equation*}
(\nabla \lambda)_{w}(x, y)=w(x, y)^{1-q}(\lambda(y)-\lambda(x)) . \tag{3}
\end{equation*}
$$

Divergence $\operatorname{div}: \mathcal{E} \rightarrow \mathcal{V}$ is:

$$
\begin{equation*}
\left(\operatorname{div}_{w} \phi\right)(x)=\frac{1}{2 d(x)^{r}} \sum_{y} w(x, y)^{q}(\phi(x, y)-\phi(y, x)) \tag{4}
\end{equation*}
$$

It is true:

$$
\langle\nabla u, \phi\rangle_{\mathcal{E}}=-\left\langle u, \operatorname{div}_{w} \phi\right\rangle_{\mathcal{V}}
$$

## Partition problem on a graph

We are interested in solving partition problems of the form

$$
\begin{equation*}
\min _{S \subset V} \sum_{(x, y) \in E: x \in V, y \in V \backslash S} w(x, y) \tag{5}
\end{equation*}
$$

## Partition problem on a graph

The problem can be expressed as

$$
\begin{equation*}
\min _{\lambda \in\{0,1\}} E^{P}(\lambda)=T V_{w}(\lambda)+\sum_{x \in V} f(\lambda(x), x), \tag{6}
\end{equation*}
$$

where

$$
\begin{gather*}
T V_{w}(\lambda)=\frac{1}{2} \sum_{x \cdot y} w(x, y)^{q}|\lambda(x)-\lambda(y)| \\
f(\lambda(x), x)=\eta(x)\left|\lambda(x)-\lambda^{0}(x)\right|^{2} \tag{7}
\end{gather*}
$$

where $\lambda^{0}$ is a binary function taking value 1 or 0 at some vertices with known classification. $\eta=0$ on unclassified vertices.

## Partition problem on a graph

Define:

$$
\begin{array}{r}
C_{s}(x)=f(0, x), \quad C_{t}(x)=f(1, x), \quad \forall x \in V \\
g(\phi(x), x)=C_{t}(x) \phi(x)+C_{s}(x)(1-\phi(x)), \quad \forall x \in V \tag{8}
\end{array}
$$

The problem

$$
\begin{equation*}
\min _{\lambda \in\{0,1\}} E^{P}(\lambda)=T V_{w}(\lambda)+\sum_{x \in V} g(\lambda(x), x) \tag{9}
\end{equation*}
$$

is equivalent to:

$$
\begin{equation*}
\min _{\lambda \in[0,1]} E^{P}(\lambda)=T V_{w}(\lambda)+\sum_{x \in V} g(\lambda(x), x) \tag{10}
\end{equation*}
$$

## Partition problem on a graph

Algorithm 1 Max-flow Algorithm
Initialize $p_{s}^{1}, p_{t}^{1}, p^{1}$ and $\lambda^{1}$. For $k=1, \ldots$ until convergence:

- Optimize $p$ flow

$$
p^{k+1}=\underset{|p(e)| \leq W(e) \forall e \in E}{\arg \max }-\frac{c}{2}\left\|\operatorname{div}_{w} p-F^{k}\right\|_{2}^{2},
$$

- Optimize source flow $p_{s}$

$$
p_{s}^{k+1}=\underset{p_{s}(x) \leq C_{s}(x) \forall x \in V}{\arg \max } \sum_{x \in V} p_{s}-\frac{c}{2}\left\|p_{s}-G^{k}\right\|_{2}^{2},
$$

- Optimize sink flow $p_{t}$

$$
p_{t}^{k+1}=\underset{p_{t}(x) \leq C_{t}(x) \forall x \in V}{\arg \max }-\frac{c}{2}\left\|p_{t}-H^{k}\right\|_{2}^{2},
$$

- Update $\lambda$

$$
\lambda^{k+1}=\lambda^{k}-c\left(\operatorname{div}_{w} p^{k+1}-p_{s}^{k+1}+p_{t}^{k+1}\right)
$$

## MNIST: test



Figure: Examples of digits from the MNIST data base
Using random initialization and random fidelity, the max-flow method obtained an accuracy of around $98.48 \%$ averaged over 100 runs with different fidelity sets of 500 randomly chosen points (or only $3.62 \%$ of the set).

## Banknote Authentication Data Set

The banknote authentication data set, from the UCI machine learning repository: http://archive.ics.uci.edu/ $\mathrm{ml} /$, is a data set of 1372 features extracted from images ( $400 \times 400$ pixels) of genuine and forged banknotes. Wavelet transform was used to extract the features from the images. The goal is to segment the banknotes into being either genuine or forged.
With the max-flow method, for a $5.1 \%$ fidelity set, we were able to obtain an average accuracy (over 100 different fidelity sets) of around $99.09 \%$.

## Two moons



## Figure: Two moons example with max-flow method

This data set is constructed from two half circles in $\mathbb{R}^{2}$ with a radius of one. The centers of the two half circles are at $(0,0)$ and $(1,0.5)$. A thousand uniformly chosen points are sampled from each circle, embedded in $\mathbb{R}^{100}$ and i.d.d. Gaussian noise with standard deviation 0.02 is added to each coordinate. Therefore, the set consists of two thousand points. Starting from some initial classification of the points, the goal is to segment the two half circles. For the max-flow method, in the case of 65 or lower number of fidelity points ( $3.25 \%$ ), we increased the number of edges of supervised points to others to avoid the trivial global minimizer where all points but the supervised ones are classified as one class.
Using random initialization and random fidelity, for the max-flow method, we obtained an average accuracy (over 100 different fidelity sets) of $97.10 \%$ and $97.05 \%$ in the case of 100 and 50 fidelity points, respectively.

## Comparison of our convex algorithms to binary MBO and GL methods

## Table: Comparison of methods

|  | max-flow | primal augmented <br> Lagrangian | binary <br> MBO | binary <br> GL |
| :---: | :---: | :---: | :---: | :---: |
| MNIST (3.6\% fidelity) random <br> initialization, random fidelity | $98.48 \%$ | $98.44 \%$ | $98.37 \%$ | $98.29 \%$ |
| MNIST (3.6\% fidelity) 2nd eigenvector <br> initialization, random fidelity | $98.48 \%$ | $98.43 \%$ | $98.36 \%$ | $98.25 \%$ |
| MNIST (3.6\% fidelity) random <br> initialization, corner fidelity | $98.47 \%$ | $98.40 \%$ | $62.35 \%$ | $64.39 \%$ |
| MNIST (3.6\% fidelity) 2nd eigenvector |  |  |  |  |
| initialization, corner fidelity | $98.46 \%$ | $98.40 \%$ | $63.87 \%$ | $63.19 \%$ |
| Banknote Data Set (5.1\% fidelity) | $99.09 \%$ | $98.75 \%$ | $95.43 \%$ | $97.76 \%$ |
| Banknote Data Set (3.6\% fidelity) | $98.83 \%$ | $98.29 \%$ | $93.48 \%$ | $96.10 \%$ |
| two moons (5\% fidelity) | $97.10 \%$ | $97.07 \%$ | $98.41 \%$ | $98.31 \%$ |
| two moons (2.5\% fidelity) | $97.05 \%$ | $96.78 \%$ | $97.53 \%$ | $98.15 \%$ |

## Test with random initialization: MNIST




Random initialization and fidelity


2nd eigenvector initialization, random


Second Eigenvector max-flow result


Second Eigenvector max-flow result


Second Eigenvector binary MBO result

## Test with random initialization: MNIST




Random initial., corner fidelity


2nd eigenvector initialization, corner

max-flow result
 max-flow result
 binary MBO result
 binary MBO result

## Tests：Rod



Random initial－ ization，random fidelity


2nd eigenvector ini－ tialization

max－flow result

binary MBO result

## Tests: Rod



Random initial., corner fidelity

max-flow result binary MBO result
 2





2nd eigenvector initialization

max-flow result
binary MBO result

Figure: Results for Rod 1. Left: initialization, supervised points are marked in yellow and green. Middle: max-flow algorithm result. Right: binary MBO result

## Number of Iterations and Timing

Table: Number of Iterations and Timing

| Number of iterations | max-flow | primal augmented <br> Lagrangian | binary MBO | binary GL |
| :---: | :---: | :---: | :---: | :---: |
| MNIST | 426 | 2709 | 10 | 52 |
| Banknote Authentication Data Set | 314 | 725 | 7 | 449 |
| two moons | 1031 | 451 | 8 | 108 |
| Timing (s) $^{\text {MNIST }}{ }^{\text {a }}$ | max-flow | primal augmented <br> Lagrangian | binary MBO | binary GL |
| Banknote Authentication Data Set | 1.21 | 43.21 | 0.52 | 0.78 |
| two moons | 4.13 | 3.76 | 0.90 | 0.95 |

${ }^{a}$ This is the timing of the method using already computed weights and eigenvalues/eigenvectors of the random walk Laplacian.

## Comparison of Final Energy

## Table: Comparison of Final Energy

| Data Set | initial energy | max-flow <br> final energy | primal augmented <br> Lagrangian <br> final energy | binary MBO <br> final energy | binary GL <br> final energy |
| :---: | :---: | :---: | :---: | :---: | :---: |
| MNIST (random fid) | 23223 | 789 | 789 | 798 | 804 |
| MNIST (non-random fid) | 23223 | 791 | 792 | 2167 | 5363 |
| Banknote Authentication | 3308 | 30 | 37 | 51 | 42 |
| two moons | 3802 | 533 | 535 | 538 | 548 |
| rod 1 (random fid) | 4159 | 146 | 148 | 163 | 159 |
| rod 1 (non-random fid) | 4159 | 88 | 89 | 825 | 391 |
| rod 2 (random fid) | 4528 | 171 | 176 | 186 | 184 |
| rod 2 (non-random fid) | 4528 | 101 | 105 | 709 | 421 |

## Cheeger ratio cut

Two-phase:

$$
\min _{\Omega \subset V} \frac{\operatorname{cut}\left(\Omega, \Omega^{c}\right)}{\min \left(\Omega, \Omega^{c}\right)}
$$

Ref: Bresson, X., Tai, X.-C. C., Chan, T. F. and Szlam, A. (2013). Multi-class Transductive Learning Based on I1 Relaxations of Cheeger Cut and Mumford-Shah-Potts Model. Journal of Mathematical Imaging and Vision, 49(1), 111.

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\min _{\Omega \subset V} \frac{\operatorname{cut}\left(\Omega, \Omega^{c}\right)}{\min \left(\Omega, \Omega^{c}\right)}
$$

Equivalent to:

$$
\min _{u \in[0,1]} \frac{T V_{w}(u)}{\|u\|_{1}} \text { s.t. } m(u):=\operatorname{median}(u)=0
$$

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$$

Equivalent to:

$$
\min _{u \in[0,1]} \max _{\lambda \in R} T V_{w}(u)-\lambda\|u\|_{1}, \quad \text { s.t. } m(u)=0
$$

Algorithm: primal-Dual.

Ref: Bresson, X., Tai, X.-C. C., Chan, T. F. and Szlam, A. (2013). Multi-class Transductive Learning Based on I1 Relaxations of Cheeger Cut and Mumford-Shah-Potts Model. Journal of Mathematical Imaging and Vision, 49(1), 111.

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$$

Equivalent to:

$$
\min _{u \in[0,1]} \max _{\lambda \in R} T V_{w}(u)-\lambda\|u\|_{1}, \quad \text { s.t. } m(u)=0
$$

Algorithm: primal-Dual.
Global convex optimization for $u$.
Ref: Bresson, X., Tai, X.-C. C., Chan, T. F. and Szlam, A. (2013). Multi-class Transductive Learning Based on I1 Relaxations of Cheeger Cut and Mumford-Shah-Potts Model. Journal of Mathematical Imaging and Vision, 49(1), 111.

## Multpiphase Cheeger ratio cut

Multiphasep:

$$
\min _{\Omega_{k} \subset V, k=1,2, \ldots K} \sum_{k} \frac{\operatorname{cut}\left(\Omega_{k}, \Omega_{k}^{c}\right)}{\min \left(\Omega_{k}, \Omega_{k}^{c}\right)}
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$$

Equivalent to:

$$
\min _{u_{k} \in[0,1]} \sum_{k} \frac{T V_{w}\left(u_{k}\right)}{\left\|u_{k}-m\left(u_{k}\right)\right\|_{1}} \text { s.t. } \sum u_{k}=1
$$

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Equivalent to:
$\min _{u_{k} \in[0,1]} \max _{\lambda_{k} \in R} \sum_{k} T V_{w}\left(u_{k}\right)-\lambda_{k}\left\|u_{k}\right\|_{1}, \quad$ s.t. $m\left(u_{k}\right)=0, \quad \sum_{k} u_{k}=1$.
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Algorithm: primal-Dual.
Global convex optimization for $u_{k}$.

## Test for 4-moon



Figure: (a) True solution, (b) Shi-Malik (c) Our algorithm.

## Road condition from online cameras



Probability：28．11\％（Dry），47．52\％（FullSnowy），Probability：9．98\％（Dry），22．08\％（FullSnowy）， 8．16\％（FullSnowylcy），2．98\％（PartialSnowylcy），26．87\％（FullSnowylcy），24．57\％（PartialSnowylcy）， 13．23\％（Wet）

16．51\％（Wet）

（h）Classification：Dry

（i）Classification：Wet
Probability：79．06\％（Dry），3．02\％（FullSnowy），Probability：4．97\％（Dry），0．05\％（FullSnowy）， 9．36\％（FullSnowylcy），4．80\％（PartialSnowylcy），0．04\％（FullSnowylcy），0．76\％（PartialSnowylcy）， 3．76\％（Wet）

94．17\％（Wet）

## Thank you!



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